

Perimeter and Coherence According to McCammond and Wise

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Abstract

In [MW97], McCammond and Wise introduce the perimeter of 2-complexes and use it to obtain a sufficient criterion for coherence. We present a new exposition of their main result, as well as some applications.

1 Introduction

A group is called *coherent* if every finitely generated subgroup is finitely presented. Coherence is a commensurability invariant, i. e., if $H < G$ is a subgroup of finite index, then G is coherent if and only if H is coherent.

Free groups and surface groups are elementary examples of coherent groups. Moreover, fundamental groups of 3-manifolds are coherent (see [Sco73]), as are mapping tori of free group automorphisms (see [FH97]).

The proofs of these statements heavily use specific properties of the groups in question. Recently, however, Wise and McCammond have developed a more general approach to coherence that applies to a wide range of examples, such as certain one-relator groups, some small-cancellation groups, etc. (see [MW97]). The first four sections of this paper contain a new exhibition of their sufficient criterion for coherence, and the remaining sections contain some applications of this criterion.

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2 Conventions and definitions

All complexes and maps between complexes are assumed to be combinatorial. $f^{(n)} \in X$ denotes an n -cell of X . Deviating from standard terminology, cells of complexes are closed, and characteristic maps extend to the boundary of a cell. In particular, for a 2-cell f of X , the domain of the characteristic map $\chi_f : D \rightarrow X$ is a polygon.

For a 2-complex X , let X_S denote its stellar subdivision. Observe that the 2-cells of X_S are triangles. For a map $\Phi : Y \rightarrow X$ between 2-complexes, let Φ_* denote the induced map at the π_1 -level and let $\Phi_S : Y_S \rightarrow X_S$ denote the induced map between the stellar subdivisions of X and Y . For an edge e of X , the *star* $St(e)$ of e is the collection of triangles of X_S that are adjacent to the edge of X_S corresponding to e .

2.1 Definition

A *weight function* on a 2-complex X is a map $w : \{\text{triangles of } X_S\} \rightarrow \mathbf{R}$. The weight $w(f)$ of a 2-cell f is the sum of the weights of the triangles in its stellar subdivision. The *standard weight function* assigns the weight 1 to every triangle in X_S . Unless stated otherwise, weight functions in this paper will take values in the nonnegative integers.

Given a weight function on X and a map $\Phi : Y \rightarrow X$, the *missing weight* of an edge e of Y is the sum of the weights of those triangles in the star of $\Phi(e)$ that are not contained in the image of the star of e . The *missing weight* of Φ is the sum of the missing weights of the edges of Y . In formulas, we have the double sum

$$M(\Phi) = \sum_{e^{(1)} \in Y} \sum_{f^{(2)} \in St(\Phi_S(e)) \setminus \Phi_S(St(e))} w(f).$$

The missing weight of a subcomplex $Y \subset X$ is the missing weight of the inclusion map.

Remark: We use the term “missing weight” instead of “perimeter” since the latter has caused some confusion.

Example: If X is a combinatorial 2-manifold without boundary (equipped with the standard weight function), then the missing weight of each 1-cell of X equals 2. If X is the 2-skeleton of the usual cubulation of \mathbf{R}^3 , then the missing weight of each 1-cell is 4, and the missing weight of each 2-cell is 12.

A 2-cell f of X with characteristic map $\chi_f : D \rightarrow X$ is said to have *exponent* n if there is a closed path w in $X^{(1)}$ such that $\chi_f|_{\partial D}$ spells the word w^n and if n is maximal. The polygon D is symmetric under rotations by $\frac{2\pi}{n}$.

The *packet* \tilde{f} of f is the complex obtained by taking n copies of D and gluing them along the boundary. There is a well defined map $\chi_{\tilde{f}} : \tilde{f} \rightarrow X$ whose restriction to the i th copy of D is χ_f precomposed with a rotation by $\frac{2\pi}{n}$.

A map $\Phi : Y \rightarrow X$ is said to be *packed* if for every 2-cell f of X and every lift γ of χ_f to Y there exists a lift of $\chi_{\tilde{f}}$ to Y extending γ . This condition is vacuous for any 2-cell whose characteristic map does not factor through Y . By attaching 2-cells to Y and extending Φ appropriately, we can assume that Φ is packed. Note that for nonnegative weights this does not increase the missing weight and that it does not change anything at the π_1 -level.

2.2 Example

Let X be the presentation complex of $\langle x \mid x^n \rangle$. X has exactly one 2-cell with characteristic map $\chi_f : D \rightarrow X$. Let $\Phi_1 : X^{(1)} \rightarrow X$ denote the inclusion of the 1-skeleton, and let Φ_2 be the restriction of χ_f to the boundary of D .

Then the missing weight of Φ_1 is exactly the weight of f , and packing Φ_1 yields a map of missing weight 0. The missing weight of Φ_2 is $n \cdot w(f)$, extending Φ_2 to the interior of D reduces the missing weight by $w(f)$, and packing Φ_2 yields a map of missing weight 0 since each copy of D reduces the missing weight by $w(f)$.

3 Reductions

Roughly speaking, the missing weight of a map $\Phi : Y \rightarrow X$ will play the role of a complexity function measuring to what extent Φ fails to be π_1 -injective.

In order to show that a finitely generated subgroup $H < \pi_1 X$ is finitely presented, we will start with some finite complex Y and a map $\Phi : Y \rightarrow X$ such that $\Phi_*(\pi_1 Y) = H$. If Φ_* has nontrivial kernel, we will attach a suitable 2-cell to Y in order to make the kernel smaller. If we can do this in such a way that the missing weight goes down with each reduction of the kernel,

this process terminates after finitely many steps, proving that H is finitely presented. If this works for any finitely generated $H < \pi_1 X$, it follows that $\pi_1 X$ is coherent, and we will give sufficient conditions for this to happen.

3.1 Definition

Fix a packed map $\Phi : Y \rightarrow X$ and a 2-cell f of X with characteristic map $\chi_f : D \rightarrow X$. Let P be an interval subdivided into no more than $|\partial D|$ edges, where $|\partial D|$ denotes the number of faces of D .

A pair $(\rho : P \rightarrow D, \rho' : P \rightarrow Y)$ of immersed paths is called a *reduction of Φ* if the following conditions are satisfied:

1. ρ and ρ' fit into the following commutative diagram:

$$\begin{array}{ccc} P & \xrightarrow{\rho'} & Y \\ \rho \downarrow & \nearrow & \downarrow \Phi \\ D & \xrightarrow{\chi_f} & X \end{array}$$

2. There is no map $D \rightarrow Y$ that will fit into the diagram above.

If $|P| < |\partial D|$, the reduction (ρ, ρ') is said to be *incomplete*, in which case we define another immersed path $\sigma : S \rightarrow \partial D$ such that S is an interval subdivided into $|S| = |\partial D| - |P|$ edges and $\rho(P) \cup \sigma(S) = \partial D$. σ is called a *complement* of ρ , and it is unique up to orientation.

If $|P| = |\partial D|$, the reduction (ρ, ρ') is called *complete*. Finally, (ρ, ρ') is said to be *maximal* if there is no reduction (τ, τ') with the property that ρ and ρ' are proper subpaths of τ and τ' , respectively.

Roughly speaking, reductions allow us to attach packets of 2-cells to Y , thus (under suitable hypotheses) reducing the missing weight of Φ . More precisely, we have the following

3.2 Construction

Fix a map $\Phi : Y \rightarrow X$ and a 2-cell f of X with characteristic map $\chi_f : D \rightarrow X$. If the pair $(\rho : P \rightarrow D, \rho' : P \rightarrow Y)$ is a reduction, then we construct a new map $\Phi^+ : Y^+ \rightarrow X$ in the following way:

If (ρ, ρ') is a complete reduction, then the endpoints of ρ are necessarily the same, even though the endpoints of ρ' in Y may not be equal. In this case, identify the endpoints of ρ' in Y , obtaining a new complex Y' . Φ factors through a map $\Phi' : Y' \rightarrow X$. Abusing notation, we still refer to the path $P \rightarrow Y \rightarrow Y'$ as ρ' . In the other case, i. e., if the endpoints of ρ' are already equal or if the reduction is incomplete, let $Y' = Y$ and $\Phi' = \Phi$. In any case, we have $M(\Phi) = M(\Phi')$.

Next, we define a complex Y'' by amalgamating D and Y' along P , i. e., $Y'' = Y' \cup_P D = (Y' \amalg D) / \sim$, where two elements $y \in Y'$ and $x \in D$ are equivalent precisely if there exists some $t \in P$ such that $\rho(t) = x$ and $\rho'(t) = y$. We can think of Y' and D as being subcomplexes of Y'' , and we define a map $\Phi'' : Y'' \rightarrow X$ by $\Phi''|_{Y'} = \Phi'$ and $\Phi''|_D = \chi_f$.

Finally, we obtain $\Phi^+ : Y^+ \rightarrow X$ by packing $\Phi'' : Y'' \rightarrow X$. Note that $\Phi^+(\pi_1 Y^+) = \Phi_*(\pi_1 Y)$.

3.3 Lemma

Let $\Phi : Y \rightarrow X$ be a packed map, and let f be a 2-cell of X with characteristic map $\chi_f : D \rightarrow X$. If $(\rho : P \rightarrow D, \rho' : P \rightarrow Y)$ is a maximal incomplete reduction, then

$$M(\Phi^+) = M(\Phi) + M(\chi_f \circ \sigma) - n \cdot w(f),$$

where n is the exponent of f , σ is a complement of ρ and Φ^+ is the map constructed in 3.2.

Proof: Choose a word w in $X^{(1)}$ such that $\chi_f|_{\partial D}$ spells the word w^n , where n is the exponent of f . If $\rho' : P \rightarrow Y$ has a closed loop as a subpath, say ρ'' , then $\Phi \circ \rho''$ necessarily spells a power of a conjugate of w in $X^{(1)}$ since $\Phi : Y \rightarrow X$ is combinatorial. This implies that the reduction ρ can be extended, which contradicts the maximality assumption. Hence, ρ' contains no closed loop.

Since ρ' contains no closed loop, the packet \tilde{f} is a subcomplex of Y^+ , and we have $M(\Phi^+|_{Y^+ \setminus \{\text{interior of } 2\text{-cells of } \tilde{f}\}}) = M(\Phi) + M(\chi_f \circ \sigma)$. Since there is no map $D \rightarrow Y$ that will fit into the diagram in definition 3.1, gluing back a 2-cell reduces the missing weight exactly by $w(f)$ (cf. example 2.2). \square

3.4 Lemma

Let $\Phi : Y \rightarrow X$ be a packed map, and let f be a 2-cell of X with characteristic map $\chi_f : D \rightarrow X$. If $(\rho : P \rightarrow D, \rho' : P \rightarrow Y)$ is a complete reduction, then

$$M(\Phi^+) \leq M(\Phi) - w(f),$$

where Φ^+ is the map constructed in 3.2.

Proof: Since ρ is complete, all the 1-cells in Y^+ correspond to 1-cells in Y , which implies that $M(\Phi) = M(\Phi^+|_{Y^+ \setminus \{\text{interior of 2-cells of } \tilde{f}\}})$. Since there is no map $D \rightarrow Y$ that will fit into the commutative diagram in definition 3.1, gluing the first 2-cell of \tilde{f} back reduces the missing weight by $w(f)$, which proves the claim. Note that (as opposed to the previous proof) ρ' may contain closed loops, in which case gluing back subsequent 2-cells may not result in a further reduction of missing weight. \square

4 A sufficient condition for coherence

Let $\alpha : Q \rightarrow X$ be a contractible immersed loop in $X^{(1)}$. Then there exists a sequence of paths $\alpha_0 = \alpha, \dots, \alpha_k = \text{const}$ such that α_{i+1} is obtained from α_i either by tightening, i. e., removing all subpaths of the form $e\bar{e}$, or by a homotopy across a 2-cell in the following way: A homotopy across a 2-cell f with characteristic map $\chi_f : D \rightarrow X$ uses a maximal reduction $(\rho : P \rightarrow D, \rho' : P \rightarrow Q)$ such that

$$\begin{array}{ccc} P & \xrightarrow{\rho'} & Q \\ \rho \downarrow & & \downarrow \alpha_i \\ D & \xrightarrow{\chi_f} & X \end{array}$$

commutes. Note that we can think of P as being a subcomplex of Q . If ρ is complete, then α_{i+1} is constructed from α_i by removing P from Q . If ρ is incomplete, then we replace the subpath $\chi_f \circ \rho$ with its complement $\chi_f \circ \sigma$ (notation as in definition 3.1; we have to choose the correct the orientation of σ in order to obtain a continuous path).

4.1 Definition

The complex X is said to have the *path reduction property* if for each contractible immersed loop α there is a sequence $\alpha_0 = \alpha, \dots, \alpha_k = \text{const}$ as above such that for every incomplete reduction (ρ, ρ') involving a 2-cell f we have the inequality

$$M(\chi_f \circ \sigma) \leq n \cdot w(f),$$

where n is the exponent of f and σ is a complement of ρ (cf. previous paragraph).

4.2 Theorem

Let X be a 2-complex enjoying the path reduction property for some weight function w . If the missing weight (with respect to w) of each edge of X is finite and if the weight of each 2-cell of X is strictly positive, then $\pi_1 X$ is coherent.

Proof: Let $H < \pi_1 X$ be a finitely generated subgroup. We want to show that H is finitely presented. There exists a finite 1-complex Y and a map $\Phi : Y \rightarrow X$ such that $\Phi_*(\pi_1 Y) = H$ and Φ is an immersion on the 1-skeleton (see [Sta83] for details).

If Φ is not π_1 -injective, then there exists an essential loop β in Y such that $\alpha = \Phi \circ \beta$ is contractible in X . Choose a sequence $\alpha_0 = \alpha, \dots, \alpha_k = \text{const}$ as above, using the path reduction property of X . The idea is to add finitely many 2-cells to Y in such a way that the loops $\alpha_0, \dots, \alpha_k$ and the homotopies between them lift to the new complex.

Assume inductively that the loops $\alpha_0, \dots, \alpha_i$ and the homotopies between them have already been lifted to $\beta_0 = \beta, \dots, \beta_i$ in Y . If α_{i+1} is obtained from α_i by tightening, then α_{i+1} lifts to Y since Φ is an immersion on the 1-skeleton.

If a homotopy across a 2-cell fails to lift to Y , then the corresponding reduction $(\rho : P \rightarrow D, \rho' : P \rightarrow Q)$ of $\alpha_i : Q \rightarrow X$ gives rise to a reduction of Φ in the following way: The path α_i in the commutative diagram

$$\begin{array}{ccc} P & \xrightarrow{\rho'} & Q \\ \rho \downarrow & & \downarrow \alpha_i \\ D & \xrightarrow{\chi_f} & X \end{array}$$

factors through Y by our inductive hypothesis, inducing a reduction $(\rho, \beta_i \circ \rho')$ of Φ :

$$\begin{array}{ccccc} P & \xrightarrow{\rho'} & Q & \xrightarrow{\beta_i} & Y \\ \rho \downarrow & & & & \downarrow \Phi \\ D & \xrightarrow{\chi_f} & & & X \end{array}$$

Extend this to a maximal reduction of Φ and form the map $\Phi^+ : Y^+ \rightarrow X$ as in 3.2. Now the loop α_{i+1} lifts to Y^+ , as does the homotopy taking α_i to α_{i+1} .

If the reduction is incomplete, then lemma 3.3 implies that $M(\Phi^+) = M(\Phi) + M(\chi_f \circ \sigma) - n \cdot w(f)$, and the path reduction property guarantees that $M(\Phi^+) \leq M(\Phi)$. If the reduction is complete, then, by lemma 3.4, $M(\Phi^+) \leq M(\Phi) - w(f) < M(\Phi)$ since $w(f) > 0$ by assumption. Φ^+ may not be an immersion on the 1-skeleton, but we can correct this by folding edges (see [Sta83]) without increasing the missing weight, which completes the inductive step.

Observe that we cannot arrive at a constant path unless at least one of the reductions of Φ is complete, so we end up with a finite complex Y' and a map $\Phi' : Y' \rightarrow X$ such that $\Phi'_*(\pi_1 Y') = H$, Φ' is an immersion on the 1-skeleton, and $M(\Phi') < M(\Phi)$.

Since we are using nonnegative integer weights, we can only repeat this process finitely many times before we arrive at a π_1 -injective map $\Phi'' : Y'' \rightarrow X$ with finite domain Y'' , which implies that H is finitely presented. Since H was arbitrary, this implies that $\pi_1 X$ is coherent. \square

5 Applications

The applications listed in this section can be found in [MW97]; we give a unified approach to them using matchings in bipartite graphs. We will use the following concepts and results graph theory (see [Wes96], sec. 3.1 for details):

5.1 Matchings in bipartite graphs

All graphs are assumed to be finite. A graph G is called *bipartite* if its vertex set V can be expressed as the disjoint union of two nonempty sets V_1, V_2 such that every edge of G connects an element of V_1 to an element of V_2 . In this case, a *matching of V_1 into V_2* is a set M of pairwise disjoint edges of G such that every vertex in V_1 is contained in some edge in M .

A bipartite graph G with bipartition V_1, V_2 is said to have the *matching property with respect to V_1* if for any subset $S \subset V_1$ the number of vertices in V_2 adjacent to S is at least the number of elements of S . The following theorem holds (see [Hal35, Wes96]):

5.2 Hall's matching condition

If G is a bipartite graph with bipartition V_1, V_2 , then G has a matching of V_1 into V_2 if and only if G has the matching property with respect to V_1 . \square

We now generalize Hall's theorem. Given a function m that assigns a positive integer to every vertex in V_1 , we call a set M of edges an *m -matching* if each element of V_2 is contained in at most one element of M , and if for every vertex v in V_1 , M contains exactly $m(v)$ edges emanating from v . We call $m(v)$ the *multiplicity* of v .

The graph G is said to have the *m -matching property* if for every subset W of V_1 , the number of elements of V_2 adjacent to elements of W is at least as large as the sum of the multiplicities of the elements of W .

5.3 The m -matching theorem

Let G be a bipartite graph with bipartition V_1, V_2 , and let m be a multiplicity function on V_1 . Then G allows an m -matching if and only if G has the m -matching property.

Proof: Let G' be a graph whose vertex set is the disjoint union of V_2 and V'_1 , where V'_1 contains vertices $v_1^x \cdots v_{m(x)}^x$ for every $x \in V_1$. For every edge (x, v) connecting $x \in V_1$ and $v \in V_2$ we choose edges $(v_1^x, v) \cdots (v_{m(x)}^x, v)$.

Now G' has the matching property if and only if G has the m -matching property, and G' has a matching of V'_1 into V_2 if and only if G has an m -matching, so the claim follows from Hall's matching theorem. \square

6 Multiplicities, matchings and coherence

Let $\mathcal{P} = \langle X \mid R \rangle$ be a finite presentation. The *incidence graph* of \mathcal{P} is the bipartite graph G with vertex set $V = R \sqcup X$ such that $r \in R$ and $x \in X$ are connected by an edge if and only if x occurs in the spelling of r (relators are assumed to be cyclically reduced). Given some multiplicity function m on R , \mathcal{P} is said to have the m -matching property if G has the m -matching property.

The following theorem is the main result of this section. It shows that, for the right choice of multiplicities, the m -matching property implies coherence. This theorem unifies the ideas behind several theorems in [MW97].

6.1 Theorem

Let the group G be given by the finite presentation \mathcal{P} satisfying the m -matching property for some multiplicity function m . If we can reduce any word representing 1 to the empty word by cyclic reduction or by replacing a subword of a relator r by its complement σ in such a way that

$$n \cdot m(r) \geq |\sigma|, \tag{1}$$

then G is coherent. Here n is the exponent of r .

Proof: Let G be the incidence graph of $\mathcal{P} = \langle X \mid R \rangle$, and let Y be the presentation complex of \mathcal{P} . By hypothesis, G has an m -matching M . Pick some edge in M , connecting some $r \in R$ to some $x \in X$. Now consider the 2-cell f of Y corresponding to r . At least one of the triangles in the stellar subdivision of f is adjacent to the 1-cell corresponding to x . Assign the weight one to one such triangle. Repeat this for all edges in M , and assign the weight zero to all remaining triangles.

Since each element of X belongs to at most one edge in M , the missing weight of each edge is at most one, so we have $|\sigma| \geq M(\sigma)$ for any path σ . Moreover, the weight of each 2-cell is exactly the multiplicity of the corresponding relator. Since inequality (1) holds, Y has the path reduction property. Since the weights of all 2-cells are positive, this shows that G is coherent. \square

When applying this theorem to a presentation \mathcal{P} , one typically begins with a multiplicity function m that satisfies (1), then one checks whether \mathcal{P}

has the m -matching property. Hence it is advantageous to choose m as small as possible. We list some reasonable choices for certain classes of presentations.

6.1.1 Dehn presentations

A presentation \mathcal{P} is said to be a *Dehn presentation* if Dehn's algorithm solves the word problem for \mathcal{P} (Dehn's algorithm solves the word problem for \mathcal{P} if any cyclically reduced word representing the identity element contains a subword u of a relator r such that $|u| > \frac{|r|}{2}$).

For a Dehn presentation \mathcal{P} and a relator r of \mathcal{P} , let n be the exponent of r and let w be a cyclically reduced word w such that $r = w^n$. We define $m(r) = \lfloor \frac{|w|-1}{2} \rfloor$ if $n = 1$, $m(r) = \lfloor \frac{|w|}{2} \rfloor$ if $n = 2$, and $m(r) = \lfloor \frac{|w|+1}{2} \rfloor$ if $n \geq 3$.

This choice of multiplicities satisfies (1) because it implies $n \cdot m(r) \geq \lfloor \frac{|r|-1}{2} \rfloor$. Hence, by theorem 6.1, a group given by a Dehn presentation is coherent if it has the matching property with respect to this choice of multiplicities.

6.1.2 Remark

Since presentations of type $C(4) - T(5)$ and $C(3) - T(7)$ define hyperbolic groups, it is natural to ask whether such presentations are Dehn presentations. The following two examples show that in general, the answer is no: Let $\mathcal{P}_1 = \langle a, b, c, d, r, s \mid ab\bar{a}r, \bar{r}bcs, \bar{s}d\bar{c}\bar{d} \rangle$ and $\mathcal{P}_2 = \langle a, b, c, d, r, s, t, u, v \mid abr, \bar{r}\bar{a}s, \bar{s}\bar{b}t, \bar{t}cu, \bar{u}dv, \bar{v}\bar{c}\bar{d} \rangle$. \mathcal{P}_1 and \mathcal{P}_2 are of type $C(4) - T(5)$ and $C(3) - T(7)$, respectively, and both \mathcal{P}_1 and \mathcal{P}_2 are presentations of the fundamental group of the closed surface of genus two, so their star graph has exactly one cycle. This implies that – up to inversion and conjugacy – there is exactly one shortest boundary label w of van Kampen diagrams with exactly one interior vertex. It is easy to check that w cannot be shortened by means of a homotopy across a relator disc, which shows that \mathcal{P}_1 and \mathcal{P}_2 are not Dehn presentations.

6.1.3 Small cancellation conditions

A presentation is said to have property P if every piece has length one and no relator is a proper power (see [GS90]).

Since a nonempty cyclically reduced word representing 1 in a $C(6)$ presentation contains a complement of no more than three pieces, inequality (1) holds if we assign the multiplicity three to all the relators in a $C(6) - P$ presentation. Hence we recover a special case of theorem 9.4 of [MW97].

A nonempty cyclically reduced word representing 1 in a $C(4) - T(4)$ presentation contains a complement of no more than two pieces, so inequality (1) holds if we assign the multiplicity two to all the relators in a $C(4) - T(4) - P$ presentation. In this case we recover a special case of theorem 9.5 of [MW97].

6.1.4 λ -presentations

A presentation is said to be a λ -presentation if every nontrivial word representing the neutral element contains a subword of some relator r of length strictly greater than $(1 - \lambda)|r|$. For example, Dehn presentations are $\frac{1}{2}$ -presentations by definition, and B.B. Newman's spelling theorem shows that presentations of the form $\langle X \mid w^n \rangle$ are $\frac{1}{n}$ -presentations if $n \geq 2$.

For a relator r with exponent n in a λ -presentation, let k be the *largest* integer satisfying $k < \lambda|r|$. Then we choose $m(r)$ to be the *smallest* positive integer satisfying $n \cdot m(r) \geq k$. By construction, inequality (1) holds.

For example, if $\mathcal{P} = \langle X \mid w^n \rangle$ for some $n \geq 2$, we have $k = |w| - 1$, and $m(r) \geq \frac{|w|-1}{n}$. The incidence graph of \mathcal{P} clearly has the m -matching property if the number of generators occurring in w is at least $\frac{|w|-1}{n}$.

In particular, the group given by \mathcal{P} is coherent if $n \geq |w| - 1$, so we recover theorem 7.3 in [MW97].

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